

Solutions - Homework 2

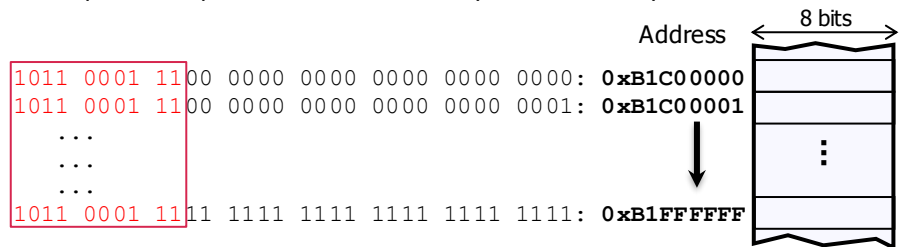
(Due date: February 6th @ 5:30 pm)

Presentation and clarity are very important! Show your procedure!

PROBLEM 1 (18 PTS)

- a) What is the minimum number of bits required to represent: (2 pts)
- ✓ 100,000 symbols? $\lceil \log_2 100,000 \rceil = 17$
 - ✓ Numbers between 0 and (including) 32678? $\lceil \log_2(32678 - 0 + 1) \rceil = 16$
- b) A microprocessor has a 32-bit address line. The size of the memory contents of each address is 8 bits. The memory space is defined as the collection of memory positions the processor can address. (6 pts)
- What is the address range (lowest to highest, in hexadecimal) of the memory space for this microprocessor? What is the size (in bytes, KB, or MB) of the memory space? $1\text{KB} = 2^{10}$ bytes, $1\text{MB} = 2^{20}$ bytes, $1\text{GB} = 2^{30}$ bytes
Address range: $0x00000000$ to $0xFFFFFFFF$.
With 32 bits, we can address 2^{32} bytes, thus we have $2^{2 \times 30} = 4$ GB of address space
 - A memory device is connected to the microprocessor. Based on the memory size, the microprocessor has assigned the addresses $0xB1C00000$ to $0xB1FFFFFF$ to this memory device.
 - What is the size (in bytes, KB, or MB) of this memory device?
 - What is the minimum number of bits required to represent the addresses only for this memory device?

As per the figure below, we only need 22 bits for the addresses in the given range. Thus, the size of the memory device is $2^{22} = 4\text{MB}$.



- c) The figure below depicts the entire memory space of a microprocessor. Each memory address occupies one byte. (10 pts)
- What is the size (in bytes, KB, or MB) of the memory space? What is the address bus size of the microprocessor?

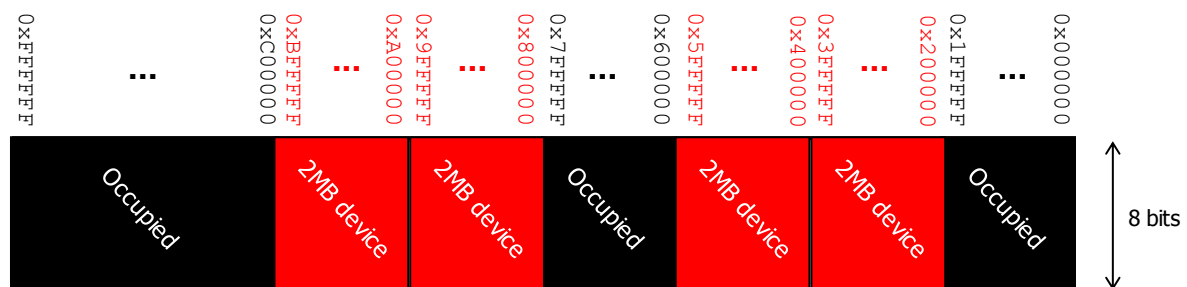
Address space: $0x000000$ to $0xFFFFFFFF$. To represent all these addresses, we require 24 bits. So, the address bus size of the microprocessor is 24 bits. The size of the memory space is then $2^{24} = 16\text{MB}$.

- If we have a memory chip of 2MB, how many bits do we require to address 2MB of memory?

2 MB memory device: $2\text{MB} = 2 \times 2^{20} = 2^{21}$ bytes. Thus, we require 21 bits to address the memory device.

- We want to connect the 2MB memory chip to the microprocessor. For optimal implementation, we must place those 2MB in an address range where every single address shares some MSBs (e.g.: $0x000000$ to $0xFFFFFFFF$). Provide a list of all the possible address ranges that the 2MB memory chip can occupy. You can only use the non-occupied portions of the memory space as shown below.

$0x200000$ to $0x3FFFFFF$
 $0x400000$ to $0x5FFFFFF$
 $0x800000$ to $0x9FFFFFF$
 $0xA00000$ to $0xBFFFFFF$



PROBLEM 4 (38 PTS)

- a) Perform the following additions and subtractions of the following unsigned integers. Use the fewest number of bits n to represent both operators. Indicate every carry (or borrow) from c_0 to c_n (or b_0 to b_n). For the addition, determine whether there is an overflow. For the subtraction, determine whether we need to keep borrowing from a higher bit. (8 pts)

Example ($n=8$):

✓ $54 + 210$

$$\begin{array}{r} 54 = 0 \times 36 = 00110110 + \\ 210 = 0 \times D2 = 11010010 \\ \hline \text{Overflow!} \rightarrow 100001000 \end{array}$$

✓ $77 - 194$

$$\begin{array}{r} 77 = 0 \times 4D = 01001101 - \\ 194 = 0 \times C2 = 11000010 \\ \hline 00001011 \end{array}$$

✓ $271 + 137$

✓ $111 + 75$

$$\begin{array}{r} \text{No Overflow } c_9=0 \\ 271 = 0 \times 10F = 100001111 + \\ 137 = 0 \times 89 = 010001001 \\ \hline 408 = 0 \times 198 = 110011000 \end{array}$$

$$\begin{array}{r} \text{Overflow! } c_9=1 \\ 111 = 0 \times 6F = 1101111 + \\ 75 = 0 \times 4B = 1001011 \\ \hline \text{Overflow!} \rightarrow 10111010 \end{array}$$

✓ $43 - 97$

✓ $128 - 43$

$$\begin{array}{r} \text{Borrow out! } b_8=1 \\ 43 = 0 \times 2B = 00101011 - \\ 97 = 0 \times 61 = 01100001 \\ \hline 0 \times CA = 11001010 \end{array}$$

$$\begin{array}{r} \text{No Borrow Out } b_8=0 \\ 128 = 0 \times 80 = 10000000 - \\ 43 = 0 \times 2B = 00101011 \\ \hline 85 = 0 \times 55 = 01010101 \end{array}$$

- b) We need to perform the following operations, where numbers are represented in 2's complement (2C): (24 pts)

✓ $413 + 617$

✓ $-97 + 256$

✓ $93 - 128$

✓ $-127 - 37$

✓ $99 - 62$

✓ $-255 - 69$

- For each case:

- ✓ Determine the minimum number of bits required to represent both summands. You might need to sign-extend one of the summands, since for proper summation, both summands must have the same number of bits.
- ✓ Perform the signed (2C) binary addition. The result must have the same number of bits as the summands.
- ✓ Determine whether there is overflow by:
 - Using c_n, c_{n-1} (carries).
 - Performing the operation in the decimal system and checking whether the result is within the allowed range for n bits, where n is the minimum number of bits for the summands.
- ✓ If we want to avoid overflow, what is the minimum number of bits required to represent both the summands and the result?

$n = 11$ bits

$$\begin{array}{r} c_{11} \oplus c_{10} = 1 \\ \text{Overflow!} \\ 413 = 00110011101 + \\ 617 = 010011101001 \\ \hline 10000000110 \end{array}$$

$413 + 617 = 1030 \notin [-2^{10}, 2^{10}-1] \rightarrow \text{overflow!}$

To avoid overflow:

$n = 12$ bits (sign-extension)

$$\begin{array}{r} c_{12} \oplus c_{11} = 0 \\ \text{No Overflow} \\ 413 = 000110011101 + \\ 617 = 001001101001 \\ \hline 1030 = 01000000110 \end{array}$$

$413 + 617 = 1030 \in [-2^{11}, 2^{11}-1] \rightarrow \text{no overflow}$

$n = 8$ bits

$$\begin{array}{r} c_8 \oplus c_7 = 1 \\ \text{Overflow!} \\ -127 = 10000001 + \\ -37 = 11011011 \\ \hline 01011100 \end{array}$$

$-127 - 37 = -164 \notin [-2^7, 2^7-1] \rightarrow \text{overflow!}$

To avoid overflow:

$n = 9$ bits (sign-extension)

$$\begin{array}{r} c_9 \oplus c_8 = 0 \\ \text{No Overflow} \\ -127 = 110000001 + \\ -37 = 111011011 \\ \hline -164 = 101011100 \end{array}$$

$-127 - 37 = -164 \in [-2^8, 2^8-1] \rightarrow \text{no overflow}$

n = 10 bits

$c_{10} \oplus c_9 = 0$
No Overflow

$c_{10} = 1$
 $c_9 = 1$
 $c_8 = 0$
 $c_7 = 0$
 $c_6 = 0$
 $c_5 = 0$
 $c_4 = 0$
 $c_3 = 0$
 $c_2 = 0$
 $c_1 = 0$
 $c_0 = 0$

$$\begin{array}{r} -97 = 1\ 1\ 1\ 0\ 0\ 1\ 1\ 1\ 1\ 1\ 1\ + \\ 256 = 0\ 1\ 0\ 0\ 0\ 0\ 0\ 0\ 0\ 0\ 0\ 0 \\ \hline 159 = 0\ 0\ 1\ 0\ 0\ 1\ 1\ 1\ 1\ 1\ 1 \end{array}$$

$-97 + 256 = 159 \in [-2^9, 2^9-1] \rightarrow$ no overflow

n = 8 bits

$c_8 \oplus c_7 = 0$
No Overflow

$c_8 = 1$
 $c_7 = 1$
 $c_6 = 0$
 $c_5 = 0$
 $c_4 = 0$
 $c_3 = 0$
 $c_2 = 1$
 $c_1 = 0$
 $c_0 = 0$

$$\begin{array}{r} 99 = 0\ 1\ 1\ 0\ 0\ 0\ 1\ 1\ + \\ -62 = 1\ 1\ 0\ 0\ 0\ 0\ 1\ 0 \\ \hline 37 = 0\ 0\ 1\ 0\ 0\ 1\ 0\ 1 \end{array}$$

$99 - 62 = 37 \in [-2^7, 2^7-1] \rightarrow$ no overflow

n = 8 bits

$c_8 \oplus c_7 = 0$
No Overflow

$c_8 = 0$
 $c_7 = 0$
 $c_6 = 0$
 $c_5 = 0$
 $c_4 = 0$
 $c_3 = 0$
 $c_2 = 0$
 $c_1 = 0$
 $c_0 = 0$

$$\begin{array}{r} 93 = 0\ 1\ 0\ 1\ 1\ 1\ 0\ 1\ + \\ -128 = 1\ 0\ 0\ 0\ 0\ 0\ 0\ 0\ 0 \\ \hline -35 = 1\ 1\ 0\ 1\ 1\ 1\ 0\ 1 \end{array}$$

$93 - 128 = -35 \in [-2^7, 2^7-1] \rightarrow$ no overflow

n = 9 bits

$c_9 \oplus c_8 = 1$
Overflow!

$c_9 = 1$
 $c_8 = 0$
 $c_7 = 0$
 $c_6 = 0$
 $c_5 = 0$
 $c_4 = 0$
 $c_3 = 0$
 $c_2 = 1$
 $c_1 = 1$
 $c_0 = 0$

$$\begin{array}{r} -255 = 1\ 0\ 0\ 0\ 0\ 0\ 0\ 0\ 1\ + \\ -69 = 1\ 1\ 0\ 1\ 1\ 1\ 0\ 1\ 1 \\ \hline 0\ 1\ 0\ 1\ 1\ 1\ 1\ 0\ 0 \end{array}$$

$-255 - 69 = -324 \notin [-2^8, 2^8-1] \rightarrow$ overflow!

To avoid overflow:

n = 10 bits (sign-extension)

$c_{10} \oplus c_9 = 0$
No Overflow

$c_{10} = 1$
 $c_9 = 1$
 $c_8 = 0$
 $c_7 = 0$
 $c_6 = 0$
 $c_5 = 0$
 $c_4 = 0$
 $c_3 = 0$
 $c_2 = 1$
 $c_1 = 1$
 $c_0 = 0$

$$\begin{array}{r} -255 = 1\ 1\ 0\ 0\ 0\ 0\ 0\ 0\ 0\ 1\ + \\ -69 = 1\ 1\ 1\ 0\ 1\ 1\ 1\ 0\ 1\ 1 \\ \hline -324 = 1\ 0\ 1\ 0\ 1\ 1\ 1\ 1\ 0\ 0 \end{array}$$

$-255 - 69 = -324 \in [-2^9, 2^9-1] \rightarrow$ no overflow

c) Get the multiplication results of the following numbers that are represented in 2's complement arithmetic with 4 bits. (6 pts)

✓ 0100×0101 , 0110×1010 , 1011×1001 .

$$\begin{array}{r} 0\ 1\ 0\ 0\ x \\ 0\ 1\ 0\ 1 \\ \hline 0\ 1\ 0\ 0\ 0 \\ 0\ 0\ 0\ 0\ 0 \\ 0\ 1\ 0\ 0\ 0 \\ 0\ 0\ 0\ 0\ 0 \\ \hline 0\ 0\ 0\ 1\ 0\ 1\ 0\ 0 \end{array}$$

$$\begin{array}{r} 1\ 0\ 0\ 0\ x \\ 0\ 1\ 1\ 0 \\ \hline 0\ 0\ 0\ 0\ 0 \\ 1\ 0\ 0\ 0\ 0 \\ 1\ 0\ 0\ 0\ 0 \\ 0\ 0\ 0\ 0\ 0 \\ \hline 0\ 0\ 1\ 1\ 0\ 0\ 0\ 0 \\ \hline 1\ 1\ 0\ 1\ 0\ 0\ 0\ 0 \end{array}$$

$$\begin{array}{r} 1\ 0\ 0\ 1\ x \\ 1\ 0\ 0\ 1 \\ \hline 0\ 1\ 1\ 1\ 1 \\ 0\ 1\ 1\ 1\ 1 \\ 0\ 1\ 1\ 1\ 1 \\ 0\ 0\ 0\ 0\ 0 \\ \hline 0\ 0\ 1\ 1\ 0\ 0\ 0\ 1 \end{array}$$